

AdKDD 2019, August 5, 2019

Feasible Bidding Strategies through Pure-Exploration Bandits

Julian Katz-Samuels, **Abraham Bagherjeiran (*)**

(*) Amazon Advertising – Sponsored Brands, Palo Alto, CA

Bidding in online advertising

A three-sided marketplace



Publisher

Requests ads, runs auctions, wants more ad spend



Advertiser

Bids for ad opportunities across multiple publishers to drive their business



User

Interacts with the ad (or not), wants meaningful ad experiences

Bidding strategies

A strategy is an algorithm that real-time bidders use to calculate the bid for an ad opportunities. Different strategies may be needed depending on the context

- First-price auctions
- Dynamic floor pricing
- Marketplace reserve pricing
- Changes to ad campaign setup



Selecting a bidding strategy

Subject of this paper

We are not proposing a new bidding strategy but a process for selecting a few among many strategies. Why can't we just test all combinations online or offline?

- Testing strategies offline introduces errors
- Test a few strategies online
- Too many strategies to test all online
- Too many parameters to test online



amazon advertising

Many possible bidding strategies

Some example bidding strategies [21]

Fixed Bid (const)



Bid the same amount
all the time

True Value (mCPC)



Bid $p(\text{Action}) * V(\text{Action})$

Linear Scaling (lin)



Scale mCPC by a
constant, such as
historical $p(\text{Action})$

Random (rand)



Bid a random value

Run it on a trickle of traffic

Typical A/B testing scenario in online systems

Modern ad servers support A/B test where a small trickle of traffic sees a new bidding behavior

- Requires ad server to implement all strategies
- Too many options means longer tests and larger errors
- Typical A/B test can run for weeks
- How do we know which is best? By which metric?

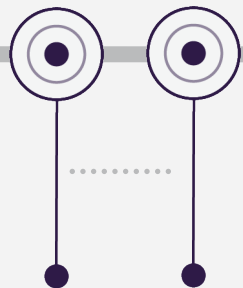


amazon advertising

Should we A/B test all parameters of bidding strategies

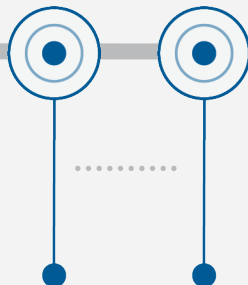
Might be nice, but there are too many parameters to test online

Fixed Bid (const)



What amount to bid

Linear (lin)



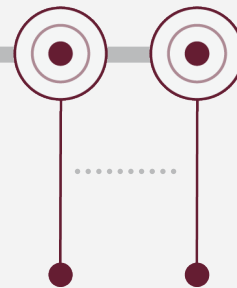
Scaling factor

True Value (mCPC)



No parameters

Random (rand)



Mean and variance

Sounds like a
job for multi-
arm bandits

Traditional multi-arm bandits

Scalar-value setting

Reward distribution not known to the learner. It is observed only after each pull

- Seek to maximize total long-term value
- Explore-exploit
- Explore more at first
- Exploit more at the end

Notation

K arms $\{1, \dots, K\}$

ν_i scalar-valued reward distribution

$$\mu_i = \mathbb{E}_{X \sim \nu_i} X$$

I_t arm pulled at each round, t

$X_t \sim \nu_{I_t}$ reward observed after each pull

Pure-exploration bandits

See [11, 15]

Explore until a fixed, predetermined stopping criteria

- Fixed confidence [1, 3, 10, 12, 14]
- Fixed budget [3, 13]
- Explore only
- No exploitation

Pure-exploration tasks

Find the one arm that maximizes μ_i after B tries

Find all arms with mean larger than a threshold, with fixed confidence

Real A/B testing, revisited

Never heard of pure-exploration bandits? Maybe you have

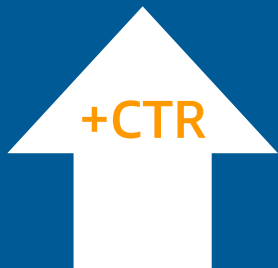
Real-world A/B testing is similar to Uniform Allocation (UA) pure-exploration bandits

- Allocate a fixed portion of traffic to each treatment
- Do not try to auto-adjust allocation (No Exploit)
- Run until B days (Fixed Budget)
- Or, run until T power (Fixed Confidence)



Acceptance criteria for A/B Tests

Determine which treatment is feasible, by meeting all criteria



Increase Click Through Rate (CTR)

Prefer bidding strategies
that surfaces ads which
will lead to clicks



Reduce Cost per Impression (CPM)

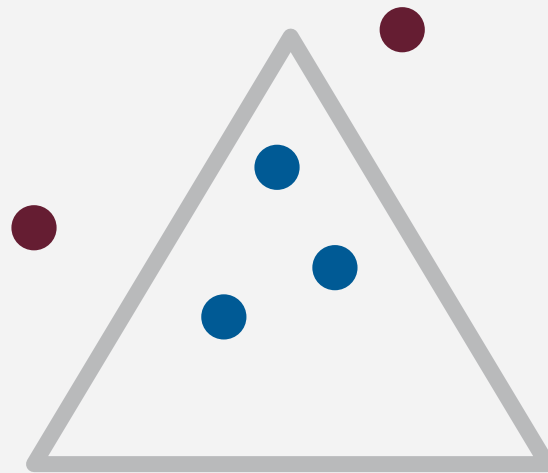
Prefer bidding strategies
that are cheaper for
advertisers

Feasible arm identification

Bandits which conform to multiple criteria

Constrained optimization for multi-armed bandits, find the best arm which is feasible.

- Multiple acceptable criteria == feasible polyhedron
- Are arms within the polyhedron or not?
- How to select arms to pull?
- Do we really need ALL the feasible arms?



Feasible polyhedron, which arm should we pull next?

Feasible arm identification

Pure exploration within fixed budget setting [13]

Now, each arm has a vector-valued reward distribution.

- Vector-valued rewards
- Feasibility == rewards within a polyhedron
- Still fixed budget
- Still explore only, no exploitation

Notation

K arms $\{1, \dots, K\}$

ν_i D -dimensional **vector**-valued
reward distribution

$$\mu_i = \mathbb{E}_{X \sim \nu_i} X$$

I_t arm pulled at each round, t

$X_t \sim \nu_{I_t}$ reward observed after
each pull

$$P = \{\mathbf{x} \in \mathbb{R}^D : A\mathbf{x} \leq b\}$$

Feasible region, polyhedron

MD-APT

Given a fixed tolerance and budget, find all feasible arms.

Pull arms which are likely to be closest to boundary

- Still uncertain about arms on the boundary
- Tolerance defines fineness of boundary
- Bounds depend on distance of all arms to boundary

Algorithm 1 MD-APT: Multi-dimensional Anytime Parameter-Free Thresholding algorithm

```
1: Input: tolerance  $\epsilon$ 
2: Initialize by pulling each arm once
3: for  $t = K + 1, \dots, T$  do
4:   Choose  $I_t = \arg \min_i [\text{dist}(\hat{\mu}_{i,t}, \partial P) + \epsilon] \sqrt{T_i(t)}$  and sample  $\mathbf{X}_t \sim \nu_{I_t}$ .
5: return  $\hat{S} = \{i \in [K] : \hat{\mu}_{i, T_i(t+1)} \in P\}$ 
```

$$\text{dist}(\hat{\mu}_{i,t}, \partial P)$$

Distance to boundary of polyhedron, P

$$\text{Define } H_\epsilon = \sum_i [\text{dist}(\mu_i, \partial P) + \epsilon]^{-2}$$

Let $K \geq 0$, $T \geq 2K$, and $\epsilon \geq 0$. Then, with probability at least

$$\Omega(1 - (\log(T) + 1)K5^D \exp(-\frac{T}{R^2 H_\epsilon})),$$

MD-APT(ϵ) returns \hat{S} such that

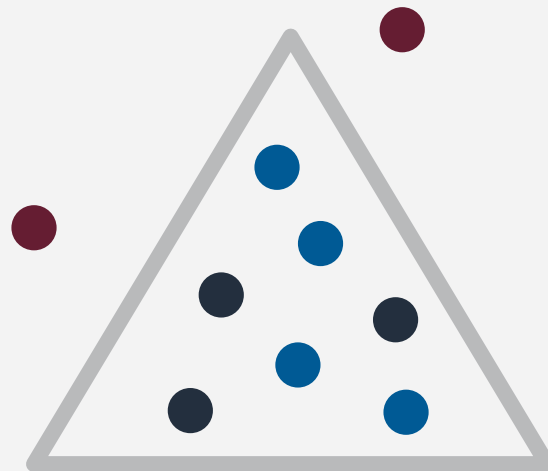
- if $\text{dist}(\mu_i, P^c) \geq \epsilon$, then $i \in \hat{S}$ and
- if $\text{dist}(\mu_i, P) \geq \epsilon$, then $i \notin \hat{S}$.

Any- m feasible arm identification

Setting discussed in this paper

Could it be faster to find only some feasible arms, but there are several edge cases

- If there are fewer than m feasible arms, return all
- Could we just end the full search earlier?



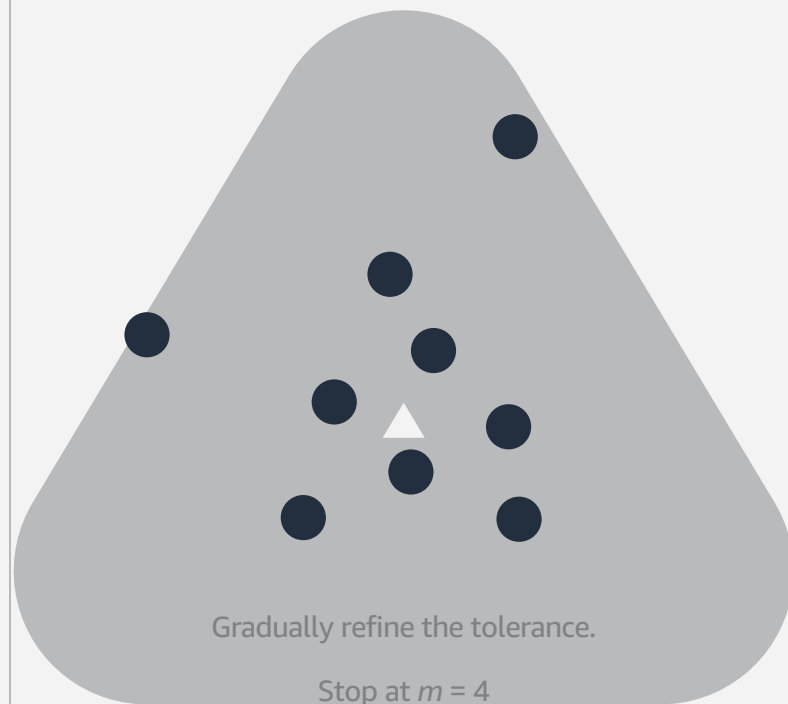
Do we really need to know
if every arm is feasible?
Maybe only $m = 4$ will do.

Any- m feasible arm identification

Example: Iteration 1

Allocate smaller budgets to MD-APT and gradually reduce tolerance

- Focus less on looking just around the boundary
- Quickly enumerates some feasible arms
- Can it guarantee returning at least m ?



Any- m feasible arm identification

Example: Iteration 2

Allocate smaller budgets to MD-APT and gradually reduce tolerance

- Focus less on looking just around the boundary
- Quickly enumerates some feasible arms
- Can it guarantee returning at least m ?



Stop at $m = 4$

Any- m feasible arm identification

Example: Iteration 3

Allocate smaller budgets to MD-APT and gradually reduce tolerance

- Focus less on looking just around the boundary
- Quickly enumerates some feasible arms
- Can it guarantee returning at least m ?



Gradually refine the tolerance.

If $m = 3$, we would be done

Any- m feasible arm identification

Example: Iteration 4

Allocate smaller budgets to MD-APT and gradually reduce tolerance

- Focus less on looking just around the boundary
- Quickly enumerates some feasible arms
- Can it guarantee returning at least m ?



Since $m = 4$, we are done

MD-APT-ANY

Given a fixed tolerance and budget, find m feasible arms.

Sequentially calls MD-APT on fixed budget to trace out arms

- Call MD-APT with decreasing tolerance
- Checks whether MD-APT found at least m -feasible arms
- Bounds depend on worst-case infeasible arm
- Returns all feasible arms if there are less than m

Algorithm 2 MD-APT-ANY

```

1: Input: tolerance  $\epsilon$ 
2: Define  $\epsilon_r := \frac{B}{2^r}$  and  $t_r = r \lceil \frac{T}{\log_2(\frac{B}{\epsilon})} \rceil$ 
3: for  $r = 1, \dots, \lceil \log_2(\frac{B}{\epsilon}) \rceil$  do
4:   Run MD-APT( $\epsilon_r$ ) for  $\frac{T}{\lceil \log_2(B) \rceil}$  iterations.
5:    $\hat{S}_r = \{i \in [K] : \hat{\mu}_{i, T_i(t_r+1)} \in P \text{ and } \text{dist}(\hat{\mu}_{i, T_i(t_r+1)}, \partial P) \geq \epsilon_{r-1}\}$ 
6:   if  $|\hat{S}_r| \geq m$  then
7:      $\hat{S} := \arg \max_{Z \subset \hat{S}_r, |Z|=m} \sum_{i \in Z} \text{dist}(\hat{\mu}_{i, T_i(t_r+1)}, \partial P)$ 
8:   Return  $\hat{S}$ 
9: Pick any  $\hat{S} \subset \{i \in [K] : \text{dist}(\hat{\mu}_{i, T_i(T+1)}, P) \leq \epsilon\}$  such that  $|\hat{S}| = \min(m, |\{i \in [K] : \text{dist}(\hat{\mu}_{i, T_i(T+1)}, P) \leq \epsilon\}|)$ .
10: Return  $\hat{S}$ 

```

$$\text{Define } H_m = \sum_{i \in [K]} \max(\Gamma_m, \text{dist}(\mu_i, \partial P))^{-2}$$

$$\Gamma_m = \max_{i: \mu_i \in P} \left(\max^{(m)} \text{dist}(\mu_i, \partial P), 0 \right)$$

Let $\epsilon > 0$. With probability at least

$$\Omega(1 - \log_2(\frac{B}{\epsilon}) \log(T) K 5^D \exp(-\frac{T}{\log_2(\frac{B}{\epsilon}) H_m R^2})),$$

MD-APT-ANY(ϵ) outputs \hat{S} such that

- if $\epsilon < \frac{\Gamma_m}{2}$, then $|\hat{S}| = m$ and $\hat{S} \subset \{i \in [K] : \mu_i \in P\}$;
- otherwise,

$$\min(|\{i \in [K] : \mu_i \in P\}|, m) \leq |\hat{S}| \text{ and } \forall i \in \hat{S} : \text{dist}(\mu_i, P) \leq 2\epsilon$$

amazon advertising

MD-APT-ANY-F

Heuristic improvement

Sample arms within border a little more on odd rounds

- Focus more on quasi-feasible arms
- Double-check that they are still feasible

Algorithm 3

1. On even rounds, run

MD-APT-ANY

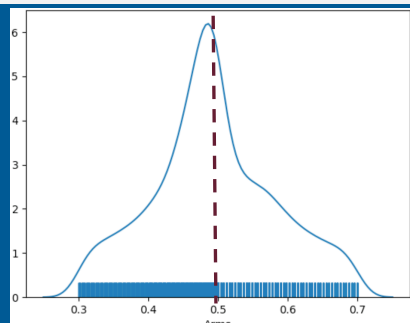
2. On odd rounds, sample the top m arms by:

$$\text{dist}(\hat{\mu}_{i, T_i(t)}, P^c) - \text{dist}(\hat{\mu}_{i, T_i(t)}, P).$$

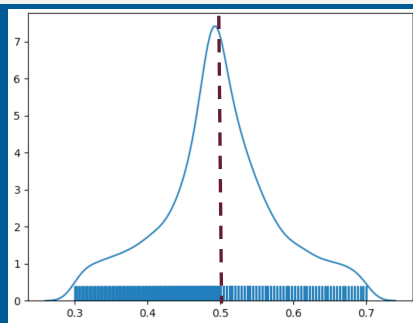
Comparing arm pull distribution

MD-APT-ANY pulls arms which are likely to be feasible more often than other algorithms

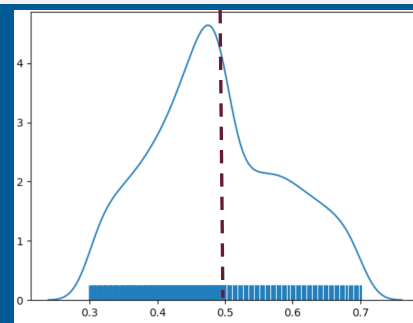
Arm pulls on synthetic experiment E2, feasible means are ≤ 0.5 (left of line below)



MD-SAR [13]



MD-APT [13]



MD-APT-ANY

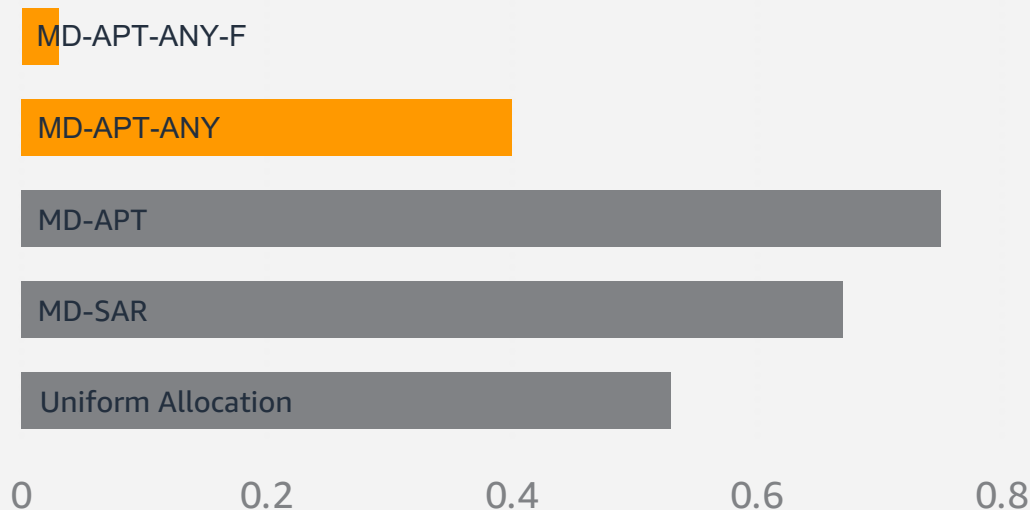
Bidding strategies data

From iPinYou RTB bidding competition [21]

- $K = 144$ arms, bidding strategies and their parameters
- $M = 1$, want to find just one good strategy
- $D = 2$, two rewards:
 - $\text{CTR} > \text{default strategy}$
 - $\text{Advertiser CPM} < \text{default strategy}$

Results on bidding strategies

Compares probability of pulling an infeasible arm



Results on other public datasets

Comparing MD-APT-ANY against other state-of-the-art feasible arm identification methods

	E1 (m = 5)	E1 (m = 10)	E1 (m = 15)	E2 (m = 5)	E2 (m = 15)	E2 (m = 20)
MD-APT-ANY-F	0.03 (0.01)	0.05 (0.01)	0.15 (0.03)	0.24 (0.03)	0.11 (0.02)	0.12 (0.02)
MD-APT-ANY	0.08 (0.02)	0.07 (0.02)	0.17 (0.03)	0.54 (0.04)	0.65 (0.03)	0.71 (0.03)
MD-APT	0.14 (0.02)	0.11 (0.02)	0.43 (0.03)	0.54 (0.04)	0.95 (0.02)	0.98 (0.01)
MD-SAR	0.14 (0.02)	0.15 (0.02)	0.34 (0.03)	0.47 (0.04)	0.80 (0.03)	0.85 (0.03)
UA	0.33 (0.03)	0.44 (0.04)	0.69 (0.03)	0.32 (0.03)	0.59 (0.03)	0.77 (0.03)

	E3 (m = 5)	E3 (m = 15)	E3 (m = 30)	Med E (m = 1)	Crowd E (m = 5)
MD-APT-ANY-F	0.20 (0.03)	0.23 (0.03)	0.35 (0.03)	0.11 (0.00)	0.18 (0.03)
MD-APT-ANY	0.29 (0.03)	0.47 (0.04)	0.36 (0.03)	0.13 (0.00)	0.32 (0.03)
MD-APT	0.36 (0.03)	0.33 (0.03)	0.40 (0.03)	0.21 (0.01)	0.62 (0.03)
MD-SAR	0.35 (0.03)	0.40 (0.03)	0.89 (0.02)	0.14 (0.00)	0.44 (0.04)
UA	0.84 (0.03)	0.82 (0.03)	0.92 (0.02)	0.13 (0.00)	0.42 (0.03)

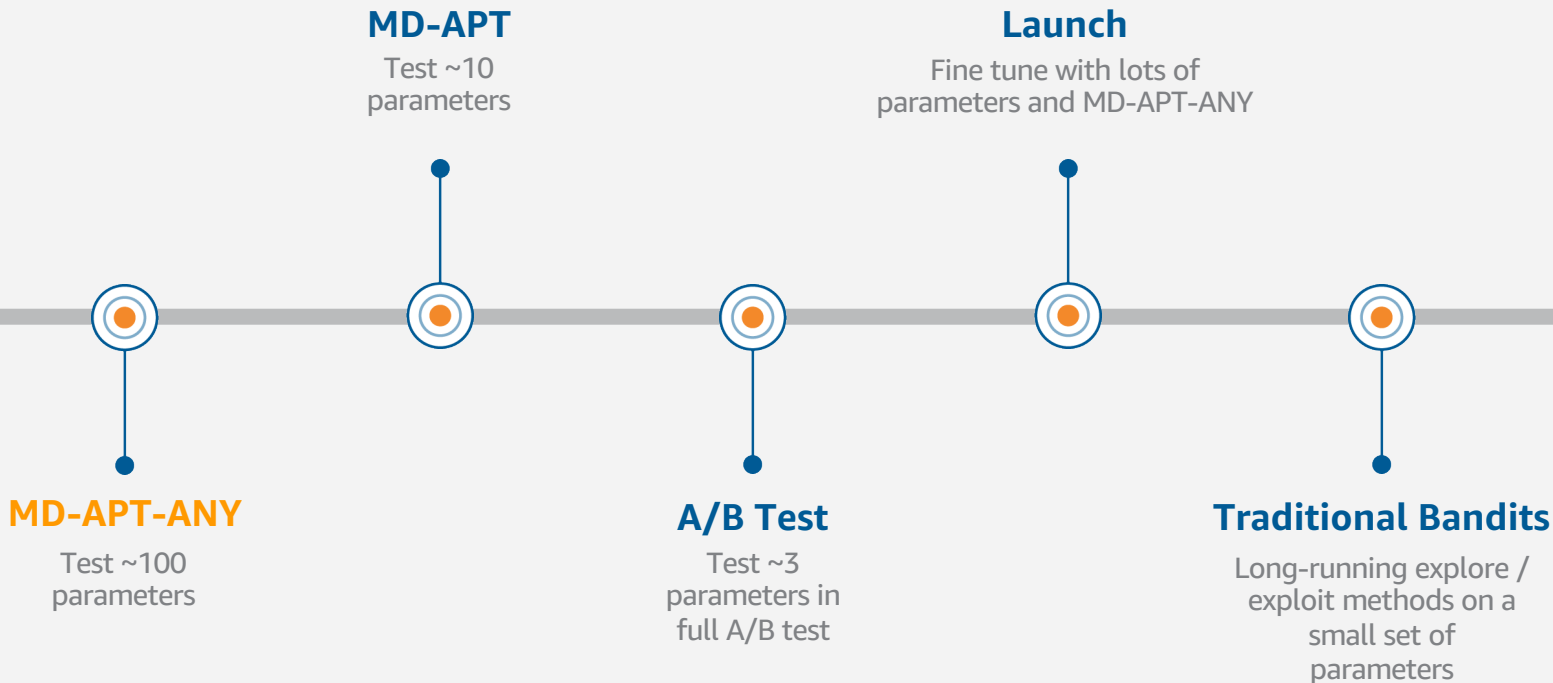
E1, E2, E3: Synthetic

Med E: Drug dosage

Crowd E: Effective of crowdsourcing workers

Hypothetical A/B testing timeline with pure-exploration bandits

Fixed-budget any-m feasible arm identification



1 Real A/B tests fit well within pure-exploration feasible-arm identification bit learning

2 Any- m feasible arm identification returns a few good parameters within budget

3 Bidding strategies and their parameters fit within this scheme, as shown in results

4 Apply to other domains such as hyperparameter optimization

Conclusion

“Join us today! We’re Hiring in Palo Alto, CA, USA”

Contact abagher@amazon.com

See our listings on Whova app

